Global Systems Science meets Programming Languages and Systems

Presentation at Second Global Systems Science Conference Stanhope Hotel, Brussels

June 11, 2013

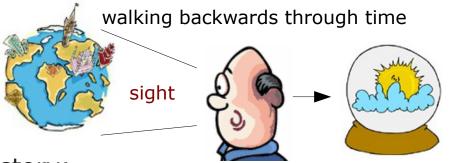
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Global Systems Science – ICT Challenges

The Past

Big data
Data grows!



Predict consequences

Big simulations

Big simulations
Increasing complexity

Complex modeling

The Future

Learn from history:

- model calibration
- analytics

You/Society **Decision Making**

ICT problems

"The free lunch is over":

Halt in CPU clock cycles

Moore's Law still holds some
years to come → more
transistors

To the rescue:

Parallel computing
Functional and declarative
 programming
Exploit domain knowledge



Some Motivation – examples

Ex 1: The Credit Crunch...

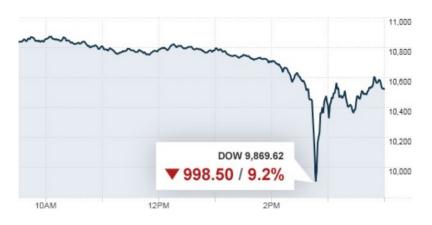


Worldwide recession in 2008: US house market collapse.

Rating ignored interdependencies and accumulated failure.

Ex 2: The Flash Crash...

Dow Jones Index on May 6, 2010:



Almost 10% drop within minutes.

Drop almost recovered minutes later.

Systemic effect of algorithmic trading at high volume and frequency.

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The HIPERFIT Research Center

Funded by the Danish Council for Strategic Research (DSF) in cooperation with financial industry partners:



2014

2015

2016

HIPERFIT: High Performance Computing for Financial IT.

Six years lifespan: 2011 2012 2013

Funding volume: 5.8M EUR.

78% funding from DSF, 22% from partners and university.

6 PhD + 3 post-doctoral positions (CS and Mathematics).

Additional funding for collaboration with small/medium-sized businesses.

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HIPERFIT Principle: "Less is More"

Transparency

Understand **more** from **shorter** code!

Understand the computation as a mathematical formula with clear semantics.

Performance

Compute more faster!

Apply domain-specific methods for parallel hardware.

Capture domain-specific parallelism in DSLs.

Productivity

Express **more** with **fewer** lines of code!

Write high-level specifications, not low-level code.

The Trick

Skip the indirection of imperative software architecture.

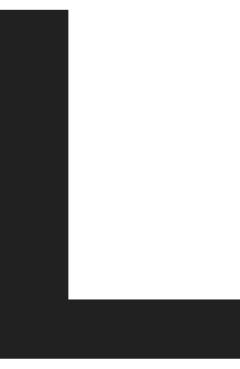
Do not build upon sequentialized inherently parallel operations!!

Use Functional Programming Language techniques!



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Vision



A High-Level, Parallel, Functional Language



Vision

Financial Contract Specification (DIKU, IMF, Nordea)

Use declarative combinators for specifying and analyzing financial contracts.

Streaming Semantics for Nested Data Parallelism (DIKU)

Reduce space complexity of "embarrassingly parallel" functional computations by streaming.

CVA (IMF, DIKU, Nordea)

Parallelize calculation of exposure to counterparty credit risk.

Automatic Parallelization of Loop Structures (DIKU)

Outperform commercial compilers on a large number of benchmarks by parallelizing and optimizing imperative loop structures.

Automatic Parallelization of Financial Applications (DIKU, LexiFi)

Analyze real-world financial kernels, such as exotic option pricing, and parallelize them to run on GPGPUs.

Bohrium (NBI)

Collect and optimize bytecode instructions at runtime and thereby efficiently execute vectorized applications independent of programming language and platform.



APL Compilation (DIKU, Insight Systems, SimCorp)

Develop techniques for compiling arrays, specifically a subset of APL, to run efficiently on GPGPUs and multicoreprocessors.

Big Data - Efficient queries (DIKU, SimCorp)

Parallelize big data queries.

Key-Ratios by Automatic Differentiation (DIKU)

Use automatic differentiation for computing sensibilities to market changes for financial contracts.

Optimal Decisions in Household Finance (IMF, Nykredit, FinE)

Investigate and develop quantitative methods to solve individual household's financial decision problems.



Project: Financial Contract Specification

Banks (and other financial institutions) use financial contracts for both

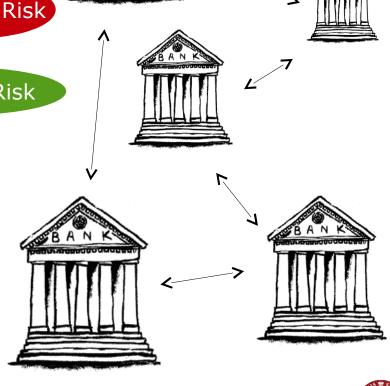
Speculation

Increase Risk

Insurance (hedging)

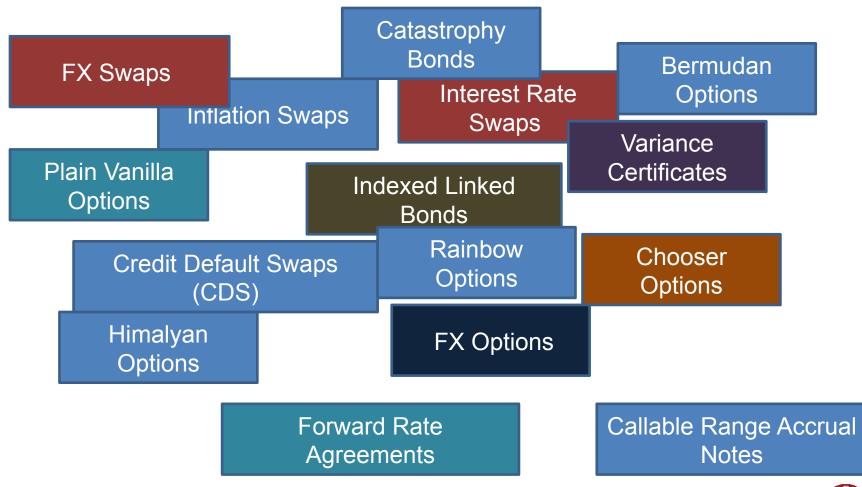
Decrease Risk

Many contracts are "Over The Counter" (OTC) contracts, which are negotiated agreements between a bank and another bank (its counter party).





Many Types of Contracts are Traded



How do Banks and Society Keep Track?

Many Problems

- Financial contracts need management *fixings, decisions, corporate actions, ...*
- Banks must **report daily** on their total value of assets
- Banks must control risk (counterparty risk, currency risk,].
- Banks need to know about future cash flows, ...

A Solution:

- Specify financial contracts in a domain specific language
- Use a functional programming language (e.g., ML)

Financial industry has already recognized the value of FP:

- LexiFi (ICFP'00 paper by Peyton-Jones, Eber, Seward)
- SimCorp A/S (uses LexiFi technology)
- Jane Street Capital (focus on electronic trading)
- Societe Generale, Credit Suisse, Standard Chartered
- Contract "Pay-off" specifications are often functional in style

Algebraic Properties, simple reasoning



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Example Financial Contracts in an Embedded DSL

```
(* Simple amortized loan *)
                                            Amortized Loan
val ex1 =
    let val coupon = 11000.0
                                             30,000 up front
        val principal = 30000.0
    in All [Give(flow(?"2011-01-01", principal, EUR)),
             flow(?"2011-02-01", coupon, EUR),
             flow(?"2011-03-01", coupon, EUR),
             flow(?"2011-04-01", coupon, EUR)]
    end
(* Cross currency swap *)
val ex2 =
    All [Give(
           All[flow(?"2011-01-01", 7000.0, DKK),
                flow(?"2011-02-01",7000.0,DKK),
                flow(?"2011-03-01",7000.0,DKK)]),
         flow(?"2011-01-01",1000.0,EUR),
         flow(?"2011-02-01",1000.0,EUR),
         flow(?"2011-03-01",1000.0,EUR)]
                                                Notice: flows in
                                              Different currencies
```

BANK

11,000 each month for 3 months

A More Complex Example...

```
(* Call option on "Carlsberg" stock *)
val equity = "Carlsberg"
val maturity = ?"2012-01-01"
val ex4 =
    let val strike = 50.0
        val nominal = 1000.0
        val obs =
            Max(Const 0.0,
                 Sub (Underlying (equity, maturity),
                     Const strike))
    in Scale (Const nominal,
             Acquire (maturity, Scale (obs, One EUR)))
    end
```

Meaning: Acquire at maturity the amount (in EUR), calculated as follows (*P* is price of Carlsberg stock at maturity): nominal * max(0, *P* – strike)



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What can we do with the contract definitions?

- Report on the expected future cash flows
- Perform management operations:
 - Advancement (simplify contract when time evolves)
 - Corporate action (stock splits, merges, catastrophic events, ...)
 - Perform fixing (simplify contract when an underlying becomes known)
- Report on the value (price) of a contract
 - Stochastic models; Monte-Carlo simulation
- Calculate risk (Key-Ratios, MC Var, CVA)



Project: Compiling APL

APL is in essence a functional language

APL has arrays as its primary data structure

APL "requires a special keyboard"!

APL is a mistake, carried through to perfection. It is the language of the future for the programming techniques of the past: it creates a new generation of coding bums.

Edsger Dijkstra

Examples:

$$a \leftarrow \iota 8$$
 $\cap array [1..8]$

$$b \leftarrow +/a$$

 $b \leftarrow +/a$ \cap sum of elements in a

$$f \leftarrow \{2+\omega \times \omega\} \cap \text{function } x^2 + 2$$

$$c \leftarrow +/f$$
 "a

 $c \leftarrow +/f$ "a \cap apply f to all elements of a and sum the elements



Map



Compiling APL - An Example

APL Code:

```
diff \leftarrow \{1 \downarrow \omega^{-1} \downarrow \omega\}

signal \leftarrow \{50 [50 \downarrow 50 \times (\text{diff } 0, \omega) \div 0.01 + \omega\}

+/ signal \(\tau \) 100000
```

Generated C Code:

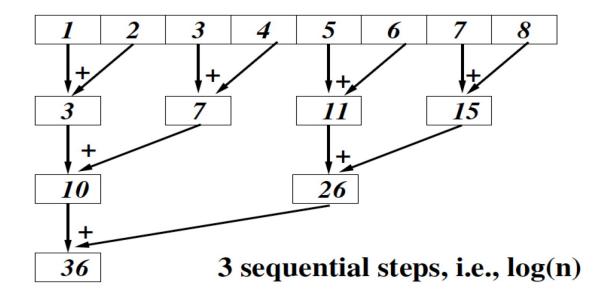
Notice: The APL Compiler has removed all notions of arrays!

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We Can Do Better!

map is "embarressingly parallel"!

Also **reduce** (+/) can be parallelized:

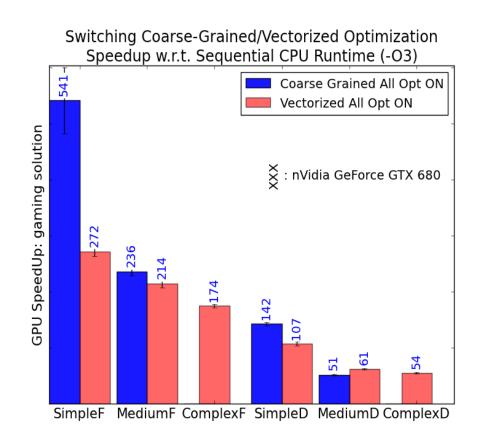


By making use of **map** and **reduce** and a few other combinators (e.g., **scan**), a high degree of parallelism can be obtained.

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Project: Pricing Financial Contracts on GPGPUs

Experiments made by HIPERFIT postdocs C. Oancea and C. Andreetta:



Three kinds of contracts: Simple, Medium, Complex

F/D: Floats/Doubles

Coarse Grained: One outer map (use of map fusion)

Vectorized: Map distribution

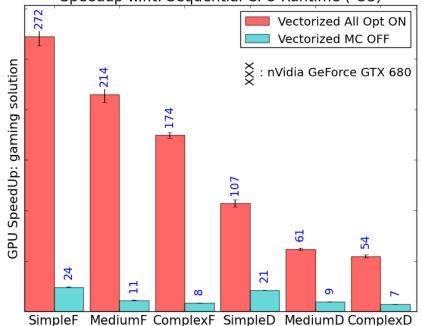
Note: Complete setup with parallel sobol sequence generation, brownian bridge, payoff function...



GPGPUs Requires New Kinds of Optimizations

Accessing global memory in a "coalesced way" may lead to dramatic speedups!

Switching ON/OFF Memory-Coalescence (MC) Optimization Speedup w.r.t. Sequential CPU Runtime (-O3)



Coalesced: consecutive threads must access consecutive global memory slots...

Often a change of algorithm is needed for ensuring coalescing (e.g., matrix transposition)...



Some Conclusions

Functional programming:

- Is declarative: Focuses on what instead of how
- Eases reasoning (formal as well as informal) and parallel processing

Open Question

The modern computation model is **highly parallel**:

- Computation everywhere simultaneously
- Grand challenge: **How to program it**?
- What is a good programming model/cost-model?

Quote (Bill Dally, senior vice president of NVIDIA Research)

"Making it easy to program a machine that requires 10 billion threads to use at full capacity is [also] a challenge. ... We need to move toward **higher-level programming models** where the programmer describes the algorithm with all available parallelism and locality exposed, and tools automate much of the process of efficiently mapping and tuning the program to a particular target machine."



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